

CS-172  
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Final

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Try to keep your answers succinct. The exam is CLOSED BOOK. All questions count equally. First, a few helpful theorems and definitions. Just because a theorem is mentioned, it may not be helpful on the exam.

**Lemma:** The *Pumping Lemma*:

If  $L$  is regular

then  $(\exists n)(\forall z \in L \text{ such that } |z| \geq n)(\exists uvw \text{ such that } z = uvw \text{ and } |uv| \leq n \text{ and } |v| \geq 1)(\forall i) : uv^i w \in L$

**Lemma:** The contrapositive of the *Pumping Lemma*:

If  $(\forall n)(\exists z \in L \text{ such that } |z| \geq n)(\forall uvw \text{ such that } z = uvw \text{ and } |uv| \leq n \text{ and } |v| \geq 1)(\exists i) : uv^i w \notin L$

then  $L$  is not regular.

**Theorem:** *Rice's theorems:* Let  $L_{\mathcal{P}}$  be the set of machines with property  $\mathcal{P}$ . If  $\mathcal{P}$  is non-trivial,  $L_{\mathcal{P}}$  is undecidable. Further,  $L_{\mathcal{P}}$  is r.e. if and only if  $\mathcal{P}$  satisfies the following three conditions:

1. If  $L \in \mathcal{P}$  and  $L \subseteq L'$  for some r.e.  $L'$ , then  $L' \in \mathcal{P}$ .
2. If  $L$  is an infinite language in  $\mathcal{P}$ , then there exists a finite subset of  $L$  in  $\mathcal{P}$ .
3. The set of finite languages in  $\mathcal{P}$  is *enumerable*.

3-SATISFIABILITY (3SAT)

**INSTANCE:** A boolean formula,  $F$ , which is an AND of clauses where each clause is an OR of 3 literals.

**QUESTION:** Is  $F$  satisfiable?

3-DIMENSIONAL MATCHING (3DM)

**INSTANCE:** A set  $M \subset W \times X \times Y$ , where  $|W| = |X| = |Y| = q$  are disjoint sets.

**QUESTION:** Does  $M$  contain a matching,  $M' \subset M$ , such that no two elements of  $M'$  agree in any coordinate.

VERTEX COVER (VC)

**INSTANCE:** A graph  $G$  and integer  $K$

**QUESTION:** Is there a subset of  $K$  vertices which cover all the edges?

CLIQUE

**INSTANCE:** A graph  $G$  and integer  $K$

**QUESTION:** Does the graph contain a clique (completely connected subgraph) of  $K$  vertices?

HAMILTONIAN CIRCUIT (HC)

**INSTANCE:** A graph  $G$

**QUESTION:** Is there a cycle through all the vertices of  $G$

PARTITION

**INSTANCE:** A finite set  $A$  and a "size"  $s(a) \in Z^+$  for each  $a \in A$ .

**QUESTION:** Is there a subset  $A' \subset A$  such that

$$\sum_{a \in A'} s(a) = \sum_{a \in A - A'} s(a)$$

1. Prove or disprove the following languages are regular:
  - (a)  $L_a = \{a^s b^t : s \geq t \geq 1\}$ .
  - (b)  $L_b = \{a^s b^t : t > s \geq 1\}$ . For the proof, use set closure properties and your result about  $L_a$ . No credit for using the pumping lemma.
  - (c)  $L_c = \{w : w \text{ contains the substring "0011"}\}$
2. Which of the following are r.e.? Give a proof. (Hint: Any reductions can be done from  $L_u$  by creating an  $M'$  from  $\langle M, w \rangle$  which accepts either  $\emptyset$  or  $\Sigma^*$  depending on whether  $M(w)$  rejects or accepts.)
  - (a)  $L_{3M} = \{\langle M_1, M_2, M_3 \rangle : \text{At least two of the machines accept the same language.}\}$
  - (b)  $\overline{L_{3M}}$
  - (c)  $L = \{\langle M \rangle : M(\epsilon) \text{ never moves past the } |Q|^{\text{th}} \text{ tape square}\}$ . ( $Q$  is the set of states of  $M$ .)
3. Of the following three problems, prove one is in NP, prove one in co-NP, and prove the third is in P.
  - (a) **INSTANCE:** Two graphs on the same vertex set  $G = (V, E)$  and  $H = (V, E')$ .  
**QUESTION:** Are  $G$  and  $H$  **non-isomorphic**?  
 (Note that it says “non-isomorphic” rather than “isomorphic”.)
  - (b) **INSTANCE:** A boolean formula,  $F$ , on the 100 variables  $\{x_1, \dots, x_{100}\}$ .  
**QUESTION:** Is  $F$  **unsatisfiable**?
  - (c) **INSTANCE:** A binary number  $n > 1$  in binary.  
**QUESTION:** Is  $n$  **composite**? (“Composite” means “not prime”).

4. Prove FEEDBACK VERTEX SET is NP-complete.

FEEDBACK VERTEX SET

**INSTANCE:** Directed graph  $G = (V, E)$  and integer  $K$ .

**QUESTION:** Is there a subset  $V' \subset V$  such that  $|V'| \leq K$  and every directed circuit in  $G$  includes at least one vertex from  $V'$ .

5. Prove HITTING STRING is NP-complete:

**INSTANCE:** An integer  $n$  and a set of strings  $A \subset \{0, 1, \#\}^n$ .

**QUESTION:** Is there a string  $x \in \{0, 1\}^n$  such that for each string  $a \in A$  there is some  $i$ ,  $1 \leq i \leq n$ , for which the  $i^{\text{th}}$  symbol of  $a$  and the  $i^{\text{th}}$  symbol of  $x$  are identical.

For example,

$$A = \{11\#0, 0\#\#\#, \#\#0\#, \#\#\#1, 0\#1\#\}$$

is a positive instance by choosing  $x = 0101$ .