
CS 70 Discrete Mathematics for CS
Fall 2001 Wagner Midterm 1

PRINT your name: _____

SIGN your name; _____

This exam is closed-book, closed-notes. One page of notes is permitted. Calculators are permitted. Do all your work on the pages of this examination.

You have 2 hours. There are 4 questions, of varying credit (50 points total). You should be able to finish all the questions, so avoid spending too long on any one question.

1. (12 pts.) Short-answer questions

Translate each of the following claims into symbolic form. For instance, a good translation of “ n is either at least three or at most five” would be “ $n \geq 3 \vee n \leq 5$.”

Then, state whether the claim is true or false, and briefly justify your answer.

(a) [3 pts.] There is some natural number whose square root is not a natural number.

(b) [4 pts.] For every natural number n , one can find another natural number m that is strictly smaller than n .

(c) [5 pts.] For each natural number k there is some lower bound ℓ so that $k^n \geq n!$ when $n \geq \ell$.

2. (12 pts.) Reachability

In chess, a bishop can move diagonally in any of the four directions. Consider a 3×3 board, with a bishop initially placed at the location marked 'B' (see below). Prove that it can never reach the square marked 'X'.

B		
		X

3. (16 pts.) Proof by induction

Let the sequence a_0, a_1, a_2, \dots be defined by the recurrence relation

$$a_n = 2a_{n-1} + a_{n-2} \text{ for } n \geq 2 \text{ and } a_0 = 1, a_1 = 2.$$

Consider the following argument:

Theorem 1 $a_n \leq n + 2$ for all $n \geq 0$.

Proof: We use strong induction on n . The base cases $n = 0$ and $n = 1$ hold, since $a_0 = 1 \leq 0 + 2$ and $a_1 = 2 \leq 1 + 2$. Now if $a_i \leq i + 2$ for each $i = 0, 1, \dots, n - 1$, for some $n \geq 2$, then we have

$$a_n = 2a_{n-1} + a_{n-2} \leq 2((n-1) + 2) + ((n-2) + 2) = 2n + 2 + n = n + 2,$$

which shows that $a_n \leq n + 2$ holds for all $n \geq 0$. \square

(a) [6 pts.] Critique the above proof.

(b) [10 pts.] Give a better proof of the theorem.

4. (10 pts.) Matchings

Recall that a *matching* on n boys and m girls is a pairing where each boy is married to exactly one girl and each girl is married to exactly one boy.

- (c) [5 pts.] Let M be a stable matching on n boys and n girls where Alice is paired with Bob. Now Alice and Bob fly off the Bermuda on vacation. We are left with a matching, call it L , on the remaining $n-1$ boys and $n-1$ girls according to who is still paired up. Is L guaranteed to be a *stable* matching, if M is stable? Prove your answer.

- (d) [5 pts.] If M, M' are two matchings, let $M \sqcap M'$ denote the configuration where each girl is married to the better of her two partners in M and M' (according to that girl's preference list). Is $M \sqcap M'$ guaranteed to be a matching? Prove your answer.
(Note that none of the matchings here are required to be stable.)

Finished! You're done; this is the last page; there are no more questions.